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Tameness and Abstract Elementary Classes

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- Give a basic overview of AECs
- Discuss tameness and its applications
- Pose some open questions

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Outline

- Give a basic overview of AECs
- Discuss tameness and its applications
- Pose some open questions

An important note: I've included some dates to give a sense of time frame, but there's some imprecision in the the mixing of publication dates and circulation of preprints dates, the latter being more common with more recent work.

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Beyond First Order Model Theory

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- 2 torsion modules
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Rather than exploring each class individually, the framework of Abstract Elementary Classes allows one to analyze them in a uniform manner.

What is an Abstract Elementary Class?

- $({\cal K},\prec_{\cal K})$ is an Abstract Elementary Class (AEC) iff
 - 0. every element of K is a L(K) structure;
 - $\bullet \prec_{\mathcal{K}} \text{ is a partial order on } K;$
 - **2** if $M \prec_{K} N$, then $M \subseteq N$;
 - **③** (K, \prec_K) respects L(K) isomorphisms;
 - if $M_0 \prec_K M_2$, $M_1 \prec_K M_2$, and $M_0 \subseteq M_1$, then $M_0 \prec_K M_1$;
 - Suppose (M_i ∈ K : i < α) is a ≺_K-increasing continuous chain, then
 - ∪_{i<α}M_i ∈ K and, for all i < α, we have M_i ≺_K ∪_{i<α}M_i; and
 if there is some N ∈ K so that, for all i < α, we have M_i ≺_K N, then we also have ∪_{i<α}M_i ≺_K N.; and
 - (Lowenheim-Skolem number) LS(K) is the minimal infinite cardinal above |L(K)| so for any M ∈ K and A ⊂ M, there is some N ≺_K M such that A ⊂ N and ||N|| = |A| + LS(K).

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Why Abstract Elementary Classes?

• The AEC axioms capture the model theoretic structure that exists *without the compactness theorem*.

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What's the point of AECs?

"Goal"

Shelah's Categoricity Conjecture: For every λ , there is μ_{λ} such that, if K is an AEC with $LS(K) = \lambda$ that is categorical in some cardinal $\geq \mu_{\lambda}$, then it is categorical in every cardinal $\geq \mu_{\lambda}$.

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Goal

To develop classification theory for AECs.

Shelah's Presentation Theorem

Everything I've said so far about AECs is semantic, but there is a syntactic description.

Theorem (Shelah's Presentation Theorem)

If K is an AEC with $LS(K) = \lambda$, then there is some $L_1 \supset L$ of size λ , an L_1 -theory T_1 , and a set Γ of quantifier free types such that

 $K = PC(T_1, \Gamma, L) := \{M_1 \upharpoonright L : M_1 \vDash T_1 \text{ and omits } \Gamma\}$

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Note that PC classes themselves are pretty poorly behaved: they fail

- the chain axioms;
- the existence of an LS number; and
- Shelah's Categoricity Conjecture
 Silver showed there is a PC class that is categorical exactly at
 κ = □_α for α limit.

Convention and embeddings

- Writing $f: M \to N$ means that f is a K-embedding, i.e. $f(M) \prec N$.
- We write K for (K, \prec_K) and \prec for \prec_K .
- We assume that K has a monster model \mathfrak{C} .
 - \mathfrak{C} is μ -model homogeneous for very large cofinality μ .
 - The existence is equivalent to amalgamation, joint embedding, and no maximal models (amalgamation is the key property most of the time)
 - Gives a very simplified definition of Galois types

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Galois types

Definition

 $A = \langle a_i : i \in I \rangle$ and $B = \langle b_i : i \in I \rangle$ have the same Galois type over M, written as gtp(A/M) = gtp(B/M), iff

there is $f \in Aut_M \mathfrak{C}$ so that $f(a_i) = b_i$ for all $i \in I$

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Definition

$$gS^{\alpha}(M) = \{gtp(\langle a_i : i < \alpha \rangle / M) : a_i \in \mathfrak{C}\}$$

If $M \prec N$ and p = gtp(a/N), then $p \upharpoonright M = gtp(a/M)$.

This definition is purely semantic. In first order, they agree with semantic types.

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Syntactict vs. Galois

- Syntactic types are very local.
- If two syntactic types differ, you can see this difference finitely: there is a finite parameter set and finite subset of the tuples that already witness the idfference

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Key Question

Do the restrictions of Galois types determine the type?

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- Give a basic overview of AECs
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As they say, a definition can't be wrong:

Definition (Grossberg-VanDieren, 2004?)

An AEC K is $< \kappa$ -tame iff for all $M \in K$ and $p, q \in gS(M)$, the two following equivalent conditions hold:

 if p ≠ q, then there is M⁻ ≺ M of size < κ such that p ↾ M⁻ ≠ q ↾ M⁻.

• if $p \upharpoonright M^- = q \upharpoonright M^-$ for all $M^- \prec M$ of size $< \kappa$, then p = q.

" κ -tameness" is "< κ ⁺-tameness."

"($< \kappa, \lambda$)-tameness" restricts the size of the domain to λ .

Tameness - A little history

- The first time something like tameness shows up is in [Sh394], where Shelah deduces *weak tameness* from categoricity and amalgamation
- Grossberg and VanDieren isolated κ-tameness in the course of the latter's thesis for an argument about Galois stability and later proved a categoricity transfer result from it
- Later authors (especially Baldwin) introduced various parameterizations and tweaks (locality, compactness, type shortness)

Variations of tameness

Tameness and locality are two of several "locality" properties for Galois types:

Definition

K is κ-local iff for all M = ∪_{i<κ}M_i and p ≠ q ∈ gS(M), there is j < κ such that p ↾ M_j ≠ q ↾ M_j.

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Tameness and locality are two of several "locality" properties for Galois types:

- K is κ-local iff for all M = ∪_{i<κ}M_i and p ≠ q ∈ gS(M), there is j < κ such that p ↾ M_j ≠ q ↾ M_j.
- K is κ-type short iff for all X, Y of size κ and M such that gtp(X/M) ≠ gtp(Y/M), there is X₀ ⊂ X and Y₀ ⊂ Y such that gtp(X₀/M) ≠ gtp(Y₀/M).

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- These can also be parameterized based on the length of types involved.
- Note I'm being vague about some of the other parameters: the length of tameness/locality/compactness and the size of the domain of type shortness

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How far from syntactic are we?

- An important/powerful class of AECs are those that are $<\kappa\text{-tame}$ and -type short for some κ
- In these classes, Galois types are determined by their restriction to small pieces, where 'small' means '< κ sized'

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- In these classes, Galois types are determined by their restriction to small pieces, where 'small' means '< κ sized'
- Doing this allows many first order arguments built on formulas to be redone in the AEC context (more on this later)
- This intuition has recently been made explicit

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Galois Morleyizations and tameness

Definition (Vasey)

Given K and κ , the $< \kappa$ -Galois Morleyization is obtained by adding predicates of lengths less than κ for all $< \kappa$ -Galois types over the emptyset.

• I can now compare the semantic $gtp_{\kappa}(a/M)$ with the syntactic $tp_{qf}(a/M^{*\kappa})$.

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Proposition (Vasey, 2015)

K is $< \kappa$ -tame and type short iff Galois types map bijectively to syntactic types in the $< \kappa$ -Galois Morleyization.

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Examples

Definition

An AEC K is $< \kappa$ -tame iff for all $M \in K$ and $p, q \in S(M)$, if $p \neq q$, then there is $M^- \prec M$ of size $< \kappa$ such that $p \upharpoonright M^- \neq q \upharpoonright M^-$.

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- quasiminimal classes are ℵ₀-tame (Zilber)
- Hrushovski fusions are ℵ₀-tame (Villaveces-Zambrano, 2005)
- Homogeneous model theory is \aleph_0 -tame
- [⊥]N is ℵ₀-tame when N is an abelian group (Baldwin-Ekloff-Trlifaj. 2007)
- torsion modules over a PID are \aleph_0 -tame (B, 2014)
- classically valued fields are ℵ0-tame (B, 2015)

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General ways of getting tameness

Proposition

If K has a "nonforking-like" notion satisfying Uniqueness, Local Character, Base Monotonicity, and Invariance, then the class is tame (for some parameters depending on what the equivalent of $\kappa(\perp)$ is).

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Theorem (B, 2013)

Let K be an AEC essentially below κ .

- If κ is weakly compact, then K is $(<\kappa,\kappa)$ -tame.
- If κ is measurable, then K is $(<\lambda,\lambda)$ -tame when cf $\lambda = \kappa$.
- If κ is nearly θ -strongly compact, then K is $(< \kappa, \theta)$ -tame.
- If κ is strongly compact, then K is $< \kappa$ -tame.

Non-examples

- For each k < ω, there is ψ_k ∈ L_{ω1,ω} that is (ℵ₀, ℵ_k)-tame, but not (ℵ_k, ℵ_{k+1})-tame. (Hart-Shelah 1990, Baldwin-Kolesnikov 2009)
- Short exact sequences of an almost free, non-free, non-Whitehead group of size κ are not (< κ, κ) tame (Baldwin-Shelah 2008)
- The large cardinals used on the previous slide are near strict (Shelah 2013?, B-Unger 2015)

Large cardinals and eventual tameness

- Global tameness principles are closely connected with large cardinal principles
- Shelah has an example showing the following:
 - If regular κ has no θ^+ -complete, uniform measure on it, there is K with $LS(K) = \theta^{\omega}$ that is not κ -local

Proposition

Suppose $\mu^{\omega} < \kappa$ for every $\mu < \kappa$.

$$\begin{pmatrix} \text{Every AEC with } LS(K) < \kappa \\ \text{is } \kappa\text{-local} \end{pmatrix} \iff \begin{pmatrix} \kappa \text{ is measurable or} \\ \text{a limit of measurables} \end{pmatrix}$$

Applications

Large cardinals and eventual tameness

Say κ is almost-strongly compact iff for every $\mu < \kappa$, every κ -complete filter can be extended to a μ -complete ultrafilter iff for all $\mu < \kappa \leq \lambda$, there is a μ -complete, fine ultrafilter on $P_{\kappa}\lambda$.

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Proposition (B-Unger, 2015)

• Suppose
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 for every $\mu < \kappa$.

$$\begin{pmatrix} \text{Every AEC with } LS(K) < \kappa \\ \text{is} < \kappa \text{-tame} \end{pmatrix} \iff \begin{pmatrix} \kappa \text{ is almost-strongly} \\ \text{compact} \end{pmatrix}$$

2

$$\left(egin{array}{c} {\sf Every} \; {\sf AEC} \; {\sf is} \ {\sf eventually} \; tame \end{array}
ight) \iff \left(egin{array}{c} {\sf There} \; {\sf are} \; class \; many \ {\sf almost-strongly} \; compact \; cardinals \end{array}
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- Pose some open questions

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Applications of tameness

There are three areas of classification theory that tameness has seen application to:

- Categoricity transfer
- Nonforking
- Stability Transfer

Applications

Categoricity Transfer

Theorem (Grossberg-VanDieren, 2006ish)

Suppose K has a monster model, is χ -tame, and categorical in some $\lambda^+ > LS(K)^+ + \chi$. Then K is categorical in all $\mu \ge \lambda^+$.

This is Shelah's Categoricity Conjecture for successors in tame AECs with a monster model.

SCC from Large Cardinals

Corollary (B, 2013)

Suppose there are class many strongly compact cardinals. If an AEC is categorical in a successor cardinal above $\mu(LS(K)) = \min\{\kappa > LS(K) : \kappa \text{ is strongly compact }\}, \text{ then it is categorical in all } \lambda \ge \mu(LS(K)).$

This uses results of Grossberg-VanDieren, Shelah, and Boney and a little more. Note that there is *no* monster model assumption.

Nonforking notions

- A main line of research is trying to find good notions of nonforking in various classes of AECs.
- Unfortunately, there's not (yet?) a single definition that specializes to all other in each circumstance.
- Still have some good results, especially when tameness holds (and a monster model exists)

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Definition

Define
$$A \stackrel{(ch)}{\underset{M_0}{\cup}} N$$
 iff for all $a \in {}^{<\kappa}A$ and $N^- \prec N$ of size $< \kappa$, $gtp(a/N^-)$ is realized in M_0 .

This is like the first-order notion of coheir, replacing "finitely satisfiable" with "< κ satisfiable."

Coheir

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Theorem (B-Grossberg, 2013)

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Suppose \kappa > LS(K). If
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- K is fully $< \kappa$ -tame and -type short;
- K does not have the κ-order property; and
 (ch)
- O \bigcirc satisfies existence/extension

then $\stackrel{(ch)}{\downarrow}$ is a "stable-like" independence relation.



If κ is strongly compact, then this is simpler.

Theorem (B-Grossberg)

Suppose $\kappa > LS(K)$ is strongly compact and $\stackrel{(ch)}{\cup}$ satisfiers existence. If K does not have the κ -order property, then κ is a "superstable-like" independence relation.

Existence (in this case) follows from categoricity λ with cf $\lambda > \kappa$.

Good λ -frames

- Shelah's focus in this area (especially recently) has been on good λ-frame s. This is a "superstable-like" notion of nonforking in a single cardinal; also comes equipped with a notion of basic types.
- Two things are done: first, prove a frame exists in some cardinal and, second, try to transfer this to larger cardinals.
- The second part uses a construction ≥ s that always exists, but doesn't always satisfy the desired properties

Both parts of this project have used non-ZFC combinatorics to get nonstructure results.

Tameness and frame existence

Previous results about the existence of frames in general required strong model and set theoretic hypotheses:

Theorem (Shelah, 2001)

lf

- $2^{\lambda} < 2^{\lambda^+} < 2^{\lambda^{++}}$ and $WDmld(\lambda^+)$ is not λ^{++} -saturated;
- K is categorical in λ and λ^+ ; and

•
$$1 \leq I(\lambda^{++}, K) < \mu_{unif}(\lambda^{++}, 2^{\lambda^+});$$

then there is a good λ^+ -frame for K.

Basic types are λ -rooted minimal types, nonforking is if the base contains the root. (Note: $\mu_{unif}(\lambda^{++}, 2^{\lambda^+})$ is "basically" $2^{\lambda^{++}}$.)

Tameness and frame existence

Tameness can replace the set-theoretic hypotheses *and* simplify the model-theoretic ones.

Theorem (Vasey, 2014)

Suppose K has a monster model, is μ -tame, and is categorical in λ with (1) cf $\lambda > \mu$ or (2) $\lambda > \mu = \beth_{\mu}$. Then K has a type-full good $\geq \lambda$ -frame.

In (1), p does not fork over M iff there is $M_0 \prec M$ of size μ so that p does not μ -split over M_0 . In (2), nonforking is μ -coheir.

Tameness and frame transfer

Previous results about the transfer of frames required strong model and set theoretic hypotheses:

Theorem (Shelah, pre-2009) If K has a good λ -frame and • $2^{\lambda} < 2^{\lambda^+} < 2^{\lambda^{++}}$ and $WDmId(\lambda^+)$ is not λ^{++} -saturated; and • $I(\lambda^{++}, K(\lambda^+ - saturated)) < \mu_{unif}(\lambda^{++}, 2^{\lambda^+});$ Then there is a good λ^+ -frame for (K', \prec') , where $K'_{\lambda^+} \subset K_{\lambda^+}$ and $\prec' \subset \prec$.

Applications

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Tameness and frame transfer

Proposition (B, 2013)

If K has amalgamation and a good λ -frame **s**, then

 \geq **s** satisfies Uniqueness iff *K* is λ -tame for basic types

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Theorem (B, B-Vasey 2014)

Suppose K has amalgamation and a good $\lambda\text{-frame}\ \mathbf{s}$ and is $\lambda\text{-tame}.$ Then

- $\mathbf{0} \geq \mathbf{s}$ is a good frame;
- (≥ s)^{<∞} is a good frame (i.e. independent sequences satisfy the nonforking properties); and
- K is $(\lambda + |\alpha|)$ -tame for basic types of length α .

More independence relations

Vasey has some recent work that looks to get global independence relations from more natural hypotheses. One version is:

Theorem (Vasey 2015)

If K has a monster model, is $< \kappa$ -tame and type short for $\kappa = \beth_{\kappa} > LS(K)$, and is categorical in $\mu > (\kappa^{<\kappa})^{+5}$, then there is a superstable-like global independence relations on models of size $\ge \mu$ and types of length $\le (\kappa^{\kappa})^{+6}$.

Stability Transfer

Applications of these concepts give rise to stability transfer results.

Theorem (Grossberg-VanDieren 2004ish)

If K is Galois stable in some $\mu > \beth_{(2^{LS(K)})^+}$ and χ -tame for $\chi < \mu$, then K is Galois stable in every $\kappa = \kappa^{\mu}$.

Theorem (Baldwin-Kueker-VanDieren 2006)

Let K with amalgamation be Galois stable in κ and κ -weakly tame. Then K is Galois stable in κ^{+n} for all $n < \omega$.

Theorem (Vasey 2014)

Suppose K is $< \chi$ -tame and stable in some $\mu \ge \chi$. Then there is some $\kappa < \beth_{(2\chi)^+}$ such that K is stable in all $\lambda \ge \mu$ such that $\lambda^{<\kappa} = \lambda$.

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- Give a basic overview of AECs
- Discuss tameness and its applications
- Pose some open questions

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- Tameness is a relatively recent notion (2004?)
- Still lots of unanswered questions and open problems

Structural Property vs. Model theoretic property

Is tameness a structural property/dividing line OR is it just a model theoretic property? That is, can we find some non-structure from non-tameness or is it just something that some AECs have and some don't.

Structural Property vs. Model theoretic property

Is tameness a structural property/dividing line OR is it just a model theoretic property? That is, can we find some non-structure from non-tameness or is it just something that some AECs have and some don't.

Vasey's result on frame existence can be rephrased as a partial answer in the good direction: Suppose K is an AEC with a monster model and is categorical in a high enough cardinal, then

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K is \mu-tame for some \mu
iff
K has a good \geq \chi-frame for some \chi
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More examples and applications

- Examples of AECs is a pretty underdeveloped field.
- Can we find more examples of non-tame AECs?
- Find some concrete (and mathematically interesting) AECs that are tame and apply the above independence relations/ideas.

Less global tameness principles

- We saw that global tameness principles were large cardinals in disguise.
- Is there any hope of getting ZFC tameness principles in "nice" classes of AECs?
 - e.g. the class can be defined recursively or in a particular descriptive set-theoretic class
- Phrased another way: All the known non-examples are pathological in one way or another. Is there a natural AEC that is not tame?

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Length of Tameness

We could parameterize tameness based on the length of tuples.

Definition

K is κ tame for α -types iff for every $p, q \in gS^{\alpha}(M)$, if $p \neq q$, there is $M^{-} \prec M$ of size κ such that $p \neq q$.

Obviously, $\alpha < \beta$ and tameness for $\beta\text{-types}$ implies tameness for $\alpha\text{-types}.$

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Question

Does tameness for α -types imply tameness for β -types? If not, is there a natural condition that causes it to?

The work of B-Vasey on independent sequences gives a partial (but unsatisfactory) answer.

Tameness

Application

Questions



Any questions?

